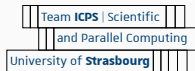


Pipelined Multithreading Generation in a Polyhedral Compiler

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Motivating Example

```
1  for (int i = 1; i < N; ++i)
2    A[i] = f1(A[i], A[i - 1]); // S1

3  for (int i = 1; i < N; ++i)
4    B[i] = f2(A[i], B[i - 1]); // S2

5  /* ... */

6  for (int i = 1; i < N; ++i)
7    F[i] = f6(E[i], F[i - 1]); // S6
```

(a) Sequential Program

Motivating Example

```
1  for (int i = 1; i < N; ++i)
2    A[i] = f1(A[i], A[i - 1]); // S1
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6
7  /* ... */
8
9  for (int i = 1; i < N; ++i)
10   F[i] = f6(E[i], F[i - 1]); // S6
```

(a) Sequential Program



(b) Dependency Graph

Motivating Example

S1(1), thread 1

```
1  for (int i = 1; i < N; ++i)
2    A[i] = f1(A[i], A[i - 1]); // S1

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(a) Sequential Program

(b) Pipelined Execution

Motivating Example



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(a) Sequential Program

(b) Pipelined Execution

Motivating Example

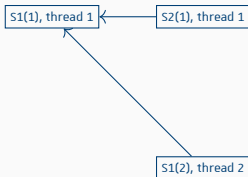
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1  for (int i = 1; i < N; ++i)
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5  /* ... */

6  for (int i = 1; i < N; ++i)
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(a) Sequential Program

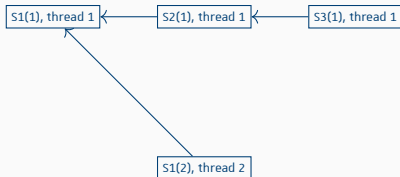


(b) Pipelined Execution

Motivating Example

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1  for (int i = 1; i < N; ++i)
2    A[i] = f1(A[i], A[i - 1]); // S1
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6  for (int i = 1; i < N; ++i)
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```

(a) Sequential Program

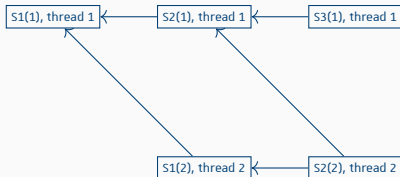


(b) Pipelined Execution

Motivating Example

```
1  for (int i = 1; i < N; ++i)
2    A[i] = f1(A[i], A[i - 1]); // S1
3  for (int i = 1; i < N; ++i)
4    B[i] = f2(A[i], B[i - 1]); // S2
5  /* ... */
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(a) Sequential Program

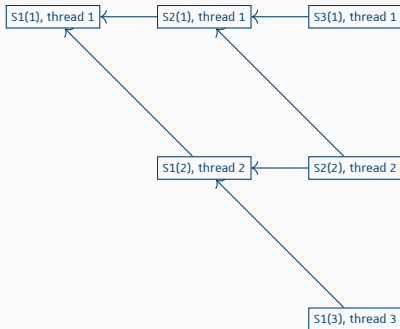


(b) Pipelined Execution

Motivating Example

```
1  for (int i = 1; i < N; ++i)
2    A[i] = f1(A[i], A[i - 1]); // S1
3  for (int i = 1; i < N; ++i)
4    B[i] = f2(A[i], B[i - 1]); // S2
5  /* ... */
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(a) Sequential Program



(b) Pipelined Execution

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1  for (int i = 1; i < N; ++i)
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7    F[i] = f6(E[i], F[i - 1]); // S6
```

(a) Sequential Program

```
1  #pragma omp parallel
2  {
3      #pragma omp for schedule(static) ordered nowait
4      for (int i = 1; i < N; ++i)
5          #pragma omp ordered
6          A[i] = f1(A[i], A[i - 1]); // S1
7      #pragma omp for schedule(static) ordered nowait
8      for (int i = 1; i < N; ++i)
9          #pragma omp ordered
10         B[i] = f2(A[i], B[i - 1]); // S2
11     /* ... */
12     #pragma omp for schedule(static) ordered nowait
13     for (int i = 1; i < N; ++i)
14         #pragma omp ordered
15         F[i] = f6(E[i], F[i - 1]); // S6
16 }
```

(b) Pipelined OpenMP target program

Motivating Example

```
1  for (int i = 1; i < N; ++i)
2    A[i] = f1(A[i], A[i - 1]); // S1
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4  for (int i = 1; i < N; ++i)
5    B[i] = f2(A[i], B[i - 1]); // S2
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7  /* ... */
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(a) Sequential Program

```
1  #pragma omp parallel
2  {
3      #pragma omp for schedule(static) ordered nowait
4      for (int i = 1; i < N; ++i)
5          #pragma omp ordered
6          A[i] = f1(A[i], A[i - 1]); // S1
7      #pragma omp for schedule(static) ordered nowait
8      for (int i = 1; i < N; ++i)
9          #pragma omp ordered
10         B[i] = f2(A[i], B[i - 1]); // S2
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12     /* ... */
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14     #pragma omp for schedule(static) ordered nowait
15     for (int i = 1; i < N; ++i)
16         #pragma omp ordered
17         F[i] = f6(E[i], F[i - 1]); // S6
18 }
```

(b) Pipelined OpenMP target program

Speedup: 2.89

6 stages on an Intel Xeon E5-2620v3 @ 2.40 GHz, with $N = 100,000$

- Identifying software pipelines in a polyhedral compiler
- Generate pipelined multithreading using OpenMP

Polyhedral Model

Introduction

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Experimental Results

Conclusion

- #pragma based API for shared memory parallelism
- Worksharing constructs
 - #pragma omp for
 - #pragma omp task

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- Worksharing constructs
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 - `#pragma omp task`
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- ♦ `#pragma` based API for shared memory parallelism
- ♦ Worksharing constructs
 - ♦ `#pragma omp for`
 - ♦ `#pragma omp task`
- ♦ Synchronization
 - ♦ `#pragma omp barrier`: explicit synchronization barrier
 - ♦ `omp_set_lock()` and `omp_unset_lock()`: explicit lock mechanism
- ♦ Clauses
 - ♦ `nowait` clause on worksharing constructs: omit the implicit barrier at the end of a worksharing construct
 - ♦ `ordered` clause on worksharing constructs: sequentialize a region

Polyhedral Model

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Sequential Loop Fission

Relaxed nowait prerequisites

Alternative: Explicit synchronization

Experimental Results

Conclusion

- Goal: maximize the number of pipeline stages
- Dependence analysis: identify Strongly Connected Components

Sequential Loop Fission

```
1  for (int i = 2; i < N; ++i) {
2    a[i] = h[i - 1] + R[i]; // S1
3    b[i] = a[i - 1] + a[i]; // S2
4    c[i] = b[i - 1] + b[i]; // S3
5    d[i] = c[i - 1] + c[i]; // S4
6    e[i] = d[i - 2] + d[i - 1]; // S5
7    f[i] = e[i - 2] + e[i - 1]; // S6
8    g[i] = f[i] + X[i]; // S7
9    h[i] = g[i] + Y[i]; // S8
10   u[i] = v[i - 1] + d[i]; // S9
11   v[i] = u[i] + Z[i]; // S10
12 }
```

(a) Original loop body

```
1  for (int i = 2; i < N; ++i) {
2    a[i] = h[i - 1] + R[i]; // S1
3    b[i] = a[i - 1] + a[i]; // S2
4    c[i] = b[i - 1] + b[i]; // S3
5    d[i] = c[i - 1] + c[i]; // S4
6    e[i] = d[i - 2] + d[i - 1]; // S5
7    f[i] = e[i - 2] + e[i - 1]; // S6
8    g[i] = f[i] + X[i]; // S7
9    h[i] = g[i] + Y[i]; // S8
10 }
11 for (int i = 2; i < N; ++i) {
12   u[i] = v[i - 1] + d[i]; // S9
13   v[i] = u[i] + Z[i]; // S10
14 }
```

(b) Fission of Strongly Connected
Components

Conditions on the `nowait` clause for parallel

The safe use of the `nowait` clause between two **parallel** loops requires that there are no dependencies between the loops or that:

- the sizes of the iteration domains are equal
- the chunk size is either the same or not specified
- both loops are bound to the same parallel region
- none of the loops is associated with a SIMD construct
- the second loop depends only on the same logical iteration of the first loop

Relaxed conditions on the `nowait` clause for ordered loops

The safe use of the `nowait` clause between two **ordered** loops requires that there are no dependencies between the loops or that:

- the sizes of the iteration domains are equal
- the chunk size is either the same or not specified
- both loops are bound to the same parallel region
- none of the loops is associated with a SIMD construct
- ~~the second loop depends only on the same logical iteration of the first loop~~

Relaxed conditions on the `nowait` clause for ordered loops

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- the sizes of the iteration domains are equal
- the chunk size is either the same or not specified
- both loops are bound to the same parallel region
- none of the loops is associated with a SIMD construct
- ~~the second loop depends only on the same logical iteration of the first loop~~
- the second loop depends on the same logical iteration **or previous logical iterations** of the first loop

Relaxed conditions on the nowait clause for ordered loops

```
1  #pragma omp parallel
2  {
3      #pragma omp for nowait
4      for (int i = 0; i < N; ++i)
5          A[i] = f1(A[i]);
6      #pragma omp for
7      for (int i = 0; i < N; ++i)
8          B[i] = f2(B[i], A[i]);
9  }
```

(a) Parallel for and nowait

```
1  #pragma omp parallel
2  {
3      #pragma omp for ordered nowait
4      for (int i = 0; i < N; ++i)
5          #pragma omp ordered
6          A[i] = f1(A[i]);
7      #pragma omp for ordered
8      for (int i = 0; i < N; ++i)
9          #pragma omp ordered
10         B[i] = f2(B[i], A[i-1]);
11 }
```

(b) Ordered for and nowait

- Annotate sequential loops with `#pragma omp for ordered`
- Enclose sequential loop bodies in `#pragma omp ordered` regions
- Annotate loops with `nowait` where possible
- Optimize by reverting `ordered` loops without `nowait` clauses to `#pragma omp single` regions

Polyhedral Model

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Sequential Loop Fission

Relaxed nowait prerequisites

Alternative: Explicit synchronization

Experimental Results

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- Loop blocking and loop fusion
- `#pragma omp for schedule(static, 1)` on the blocking loop
- `omp_set_lock()` and `omp_unset_lock()` before and after each loop of the pipeline
- up to $n \times m$ locks required for m pipeline stages over n threads

Explicit synchronization

```
1  #pragma omp parallel
2  {
3      #pragma omp for ordered nowait
4      for (size_t i = 1; i < N; ++i)
5          #pragma omp ordered
6          A[i] = f(A[i], A[i - 1]);
7
8      /* Other stages */
9
10     #pragma omp for ordered
11     for (size_t i = 1; i < N; ++i)
12         #pragma omp ordered
13         F[i] = f(E[i], F[i - 1]);
14 }
```

(a) Original program

```
1  omp_lock_t** locks;
2  #pragma omp parallel
3  {
4      /* Choose num_threads, block_size, block_count. */
5      /* Allocate, initialize and set the locks. */
6      #pragma omp for schedule(static, 1)
7      for (size_t block = 0; block < block_count; ++block)
8          /* Local loop bounds and indexes. */
9          const size_t start = 1 + block * block_size;
10         const size_t end = MIN(start + block_size, N);
11         const size_t self = block % num_threads;
12         const size_t next = (block + 1) % num_threads;
13
14         omp_set_lock(&locks[self][0]);
15         for (size_t i = start; i < end; ++i)
16             A[i] = f(A[i], A[i-1]);
17         omp_unset_lock(&locks[next][0]);
18         /* Other stages of the pipeline */
19         omp_set_lock(&locks[self][5]);
20         for (size_t i = start; i < end; ++i)
21             F[i] = f(E[i], F[i-1]);
22         omp_unset_lock(&locks[next][5]);
23     }
24 }
```

(b) Pipelined OpenMP target program

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- Tested on an Intel Xeon E5-2620v3 @ 2.40 GHz, linux 5.3.7
- Code compiled using `gcc 9.2.1` and `clang 9.0.0` with options `-O3 -march=native -fopenmp`
- FIFO scheduling enabled and process priority set to 75

benchmark	parallel loops	stages
teaser	0	5
van_dongen ¹	0	2
wdf ²	0	2
mix	1	2

¹ (Vincent H Van Dongen, Guang R Gao, and Qi Ning. “A polynomial time method for optimal software pipelining”. In: *Parallel Processing: CONPAR 92—VAPP V*. Springer, 1992, pp. 613–624)

² (Alfred Fettweis. “Wave digital filters: Theory and practice”. In: *Proceedings of the IEEE* 74.2 [1986], pp. 270–327)

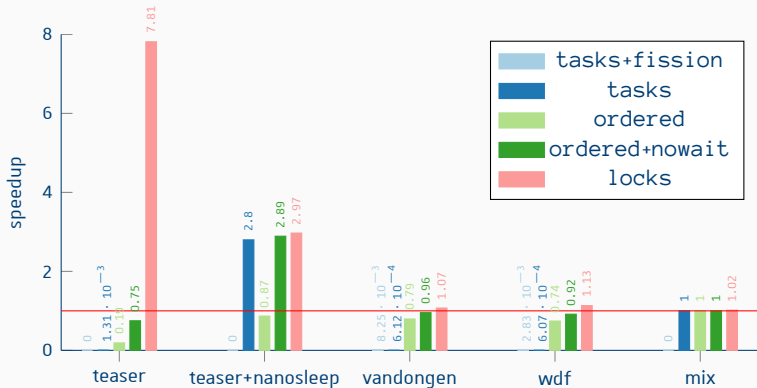


Figure 7: Speedups or slowdowns over sequential version

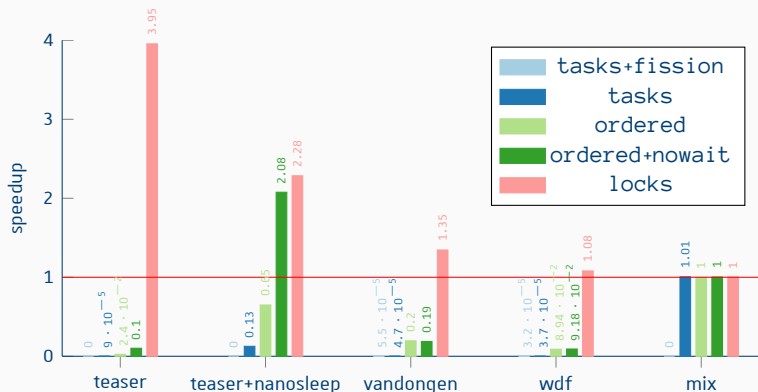


Figure 8: Speedups or slowdowns over sequential version

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- ♦ Contributions:
 - ♦ Identifying software pipelines in a polyhedral compiler
 - ♦ Generating pipelined multithreading

- ♦ Contributions:
 - ♦ Identifying software pipelines in a polyhedral compiler
 - ♦ Generating pipelined multithreading
- ♦ Future work:
 - ♦ Integration in an automatic parallelizer
 - ♦ Investigate methods to determine optimal block sizes

Appendix

- [1] Alfred Fettweis. “Wave digital filters: Theory and practice”. In: *Proceedings of the IEEE* 74.2 (1986), pp. 270–327.
- [2] Harenome Razanajato, Cédric Bastoul, and Vincent Loechner. “Pipelined Multithreading Generation in a Polyhedral Compiler”. In: *IMPACT 2020, 10th International Workshop on Polyhedral Compilation Techniques*. Bologna, Italy, Jan. 2020.
- [3] Vincent H Van Dongen, Guang R Gao, and Qi Ning. “A polynomial time method for optimal software pipelining”. In: *Parallel Processing: CONPAR 92–VAPP V*. Springer, 1992, pp. 613–624.