

Polyhedral Binary Decision Diagrams for Representing Non-Convex Polyhedra

Shubhang Kulkarni (University of Illinois, Urbana-Champaign) and Michael Kruse (Argonne National Labs)

Polyhedral Model

Simple Program

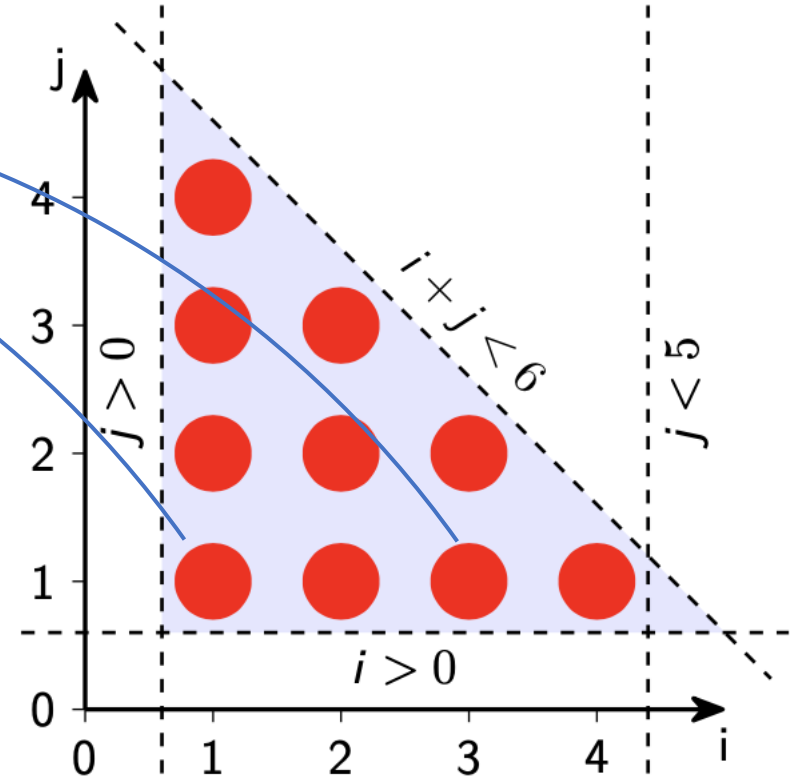
```
for (int i=1; i<5; i++)  
  for (int j=1; i+j<6; j++)  
    S(i,j);
```

Statement Domain (Explicit Enumeration)

$S(1,1), S(1,2), S(1,3), S(1,4), S(2,1), S(2,2), S(2,3),$
 $S(3,1), S(3,2), S(4,1)$

Statement Domain (Compact Representation)

$\{S(i,j) \mid 0 < i,j \wedge i+j < 6\}$



Statement Domain Visualization

Polyhedral Compilation

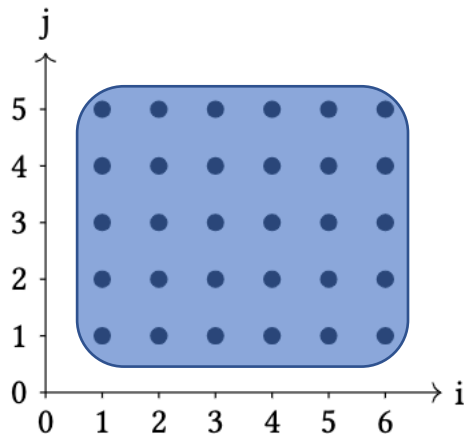
Use the polyhedral model for:

- ❖ Program Analysis
 - ❖ Check the legality of a program transform
 - ❖ E.g. Loop interchange
 - ❖ Eventually need to solve a system of linear/affine inequalities
- ❖ Statement Scheduling
 - ❖ Form a "better" schedule according to some objective
 - ❖ E.g. parallelism, data locality, memory access
 - ❖ Eventually need to solve an Integer Linear Program (IntLP)

Convexity, A Crucial Requirement

- ❖ Statement domain must be convex for compact polyhedral representation
- ❖ Does not hold for programs with conditional statements

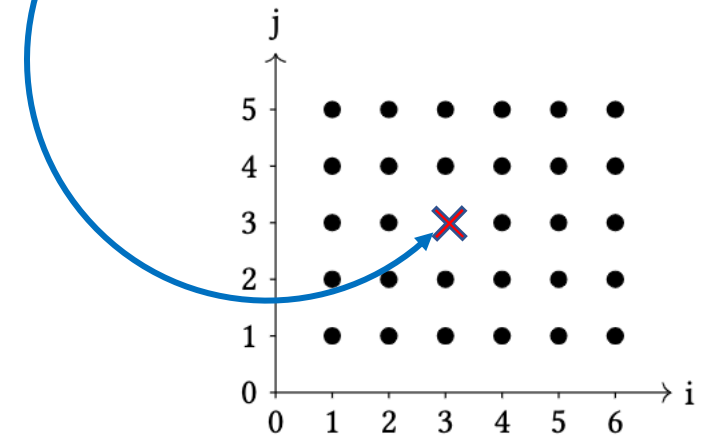
```
for (int i = 1; i <= 6; ++i)
  for (int j = 1; j <= 5; ++j)
    Stmt(i,j);
```



Programs

```
for (int i = 1; i <= 6; ++i)
  for (int j = 1; j <= 5; ++j)
    if (i != 3 || j != 3)
      Stmt(i,j);
```

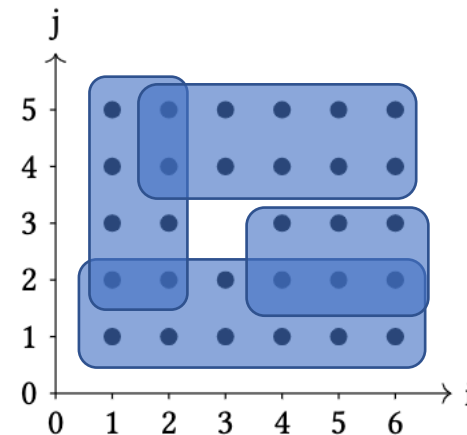
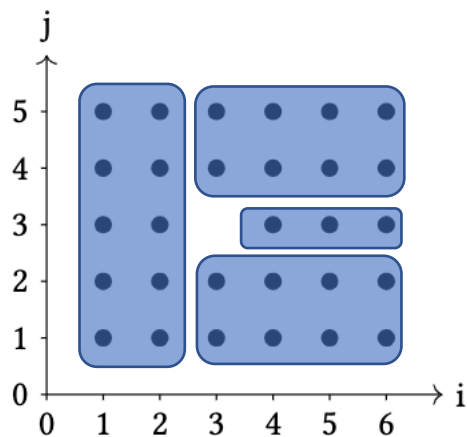
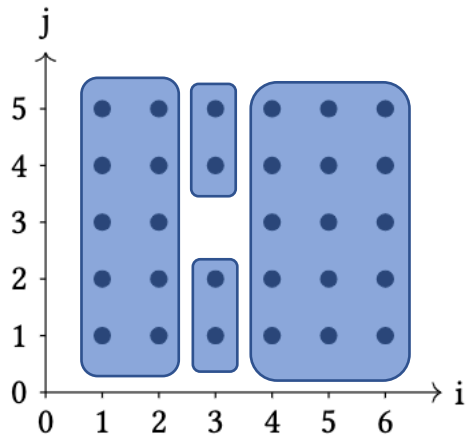
Statement Domains



Union-of-Convex-Polyhedra

❖ Non-convex vector sets are represented as union of convex polyhedra

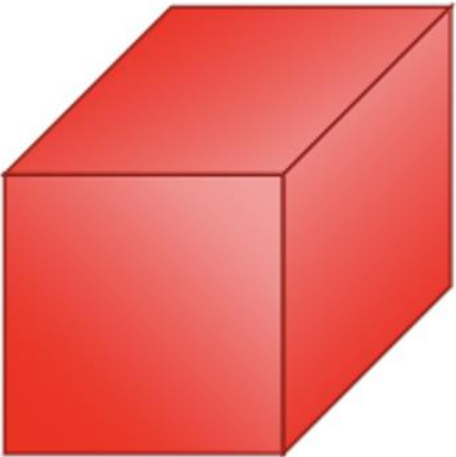
```
for (int i = 1; i <= 6; ++i)
  for (int j = 1; j <= 5; ++j)
    if (i != 3 || j != 3)
      Stmt(i,j);
```



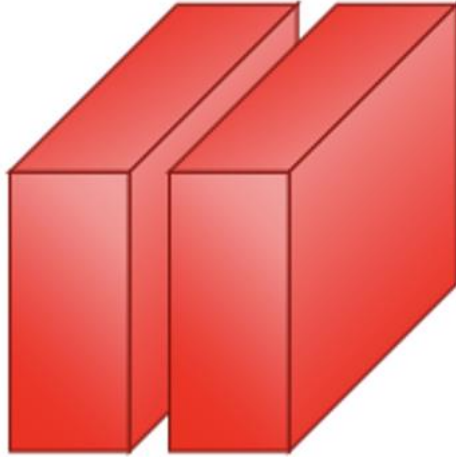
Motivating Example

❖ Conditionals introduce non-affine constraints (eg non-equality), which leads to splitting of iteration domain

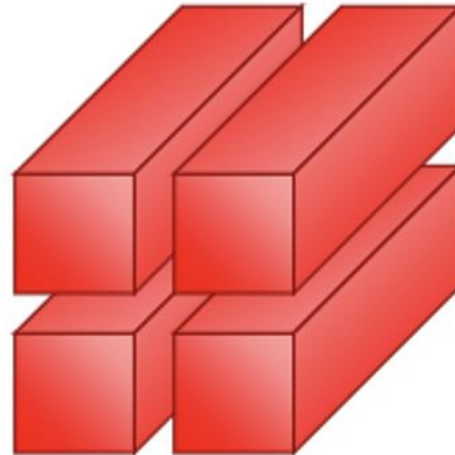
```
for (int i = 0; i < n; i+=1) {  
  for (int j = 0; j < n; j+=1) {  
    for (int k = 0; k < n; k+=1) {  
      Stmt(i, j, k);  
    }  
  }  
}
```



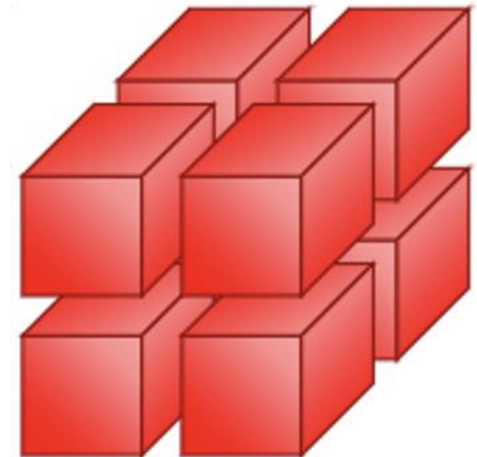
```
for (int i = 0; i < n; i+=1) {  
  for (int j = 0; j < n; j+=1) {  
    for (int k = 0; k < n; k+=1) {  
      if (i == p0)  
        continue;  
      Stmt(i, j, k);  
    }  
  }  
}
```



```
for (int i = 0; i < n; i+=1) {  
  for (int j = 0; j < n; j+=1) {  
    for (int k = 0; k < n; k+=1) {  
      if (i == p0)  
        continue;  
      if (j == p1)  
        continue;  
      Stmt(i, j, k);  
    }  
  }  
}
```



```
for (int i = 0; i < n; i+=1) {  
  for (int j = 0; j < n; j+=1) {  
    for (int k = 0; k < n; k+=1) {  
      if (i == p0)  
        continue;  
      if (j == p1)  
        continue;  
      if (k == p2)  
        continue;  
      Stmt(i, j, k);  
    }  
  }  
}
```



Polly

- ❖ The previous example shows that with p conditional statements, the union-of-convex-polyhedra representation requires $\Omega(2^p)$ polyhedra!
- ❖ Infeasible to even write down for polyhedral optimizers even when $p = 20$
- ❖ Polly is LLVM's polyhedral optimizer
- ❖ A typical execution of Polly involves
 - ❖ Representing programs via polyhedral model
 - ❖ Performing several set operations
 - ❖ Solving IntLPs

Integer Set Library (ISL)

- ❖ Polly uses Integer Set Library (ISL) to handle polyhedral computations
- ❖ ISL uses union-of-convex-polyhedral representation

Problem: Polly terminates without program-optimizations when programs have conditional-structure similar to that of motivating example.

Question: Is there an alternative representation that avoids this blowup?

YES!

Our Results

- ❖ Polyhedral Binary Decision Diagrams (PBDDs) as an alternative representation for non-convex vector sets.
- ❖ Proof of concept implementation (Python)
- ❖ Case studies comparing scaling behavior of common set operations (needed by Polly) for PBDDs vs ISL's union-of-convex-polyhedra representation

Our Results – Remarks

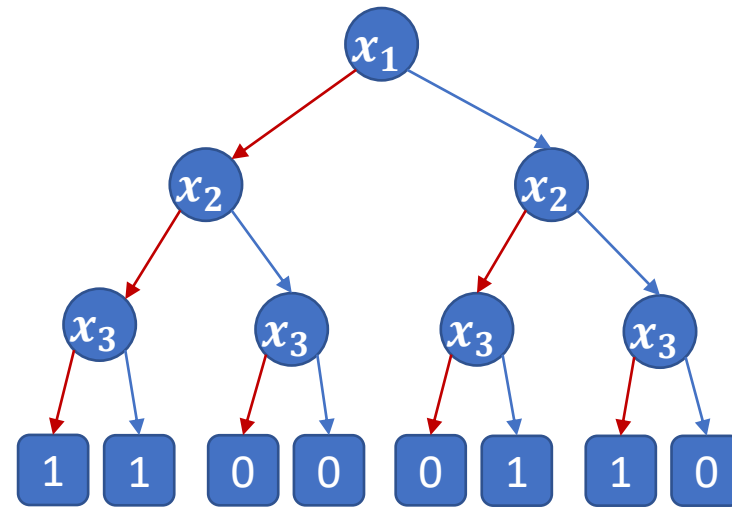
- ❖ Polyhedral Binary Decision Diagrams (PBDDs) as an alternative representation for non-convex polyhedral
 - ❖ We introduce PBDDs to make representing and performing set operations on non-convex sets feasible
 - ❖ Motivating example has a linear size PBDD representation
 - ❖ Several set-operations become computationally/algorithmically simple
 - ❖ Particularly useful when these operations result in simple sets, but Polly currently terminates prematurely due to representational overhead.
 - ❖ **Open Question:** Adapt PBDDs to work with IntLP solvers.

Binary Decision Diagrams (BDDs)

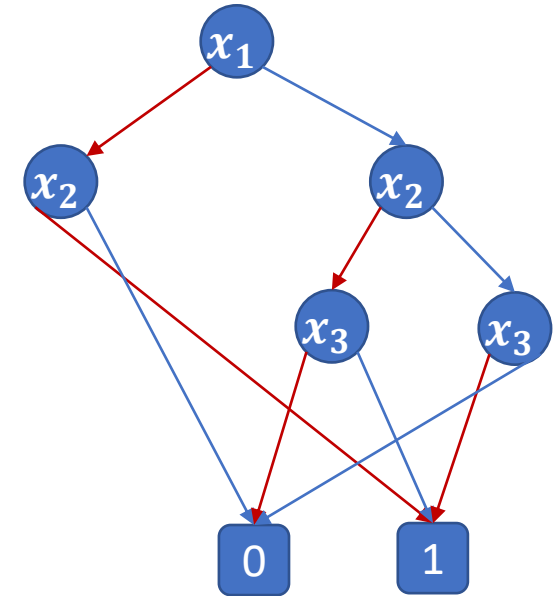
❖ Binary decision diagrams (BDDs) are used to represent Boolean functions

x_1	x_2	x_3	$f(x_1, x_2, x_3)$
0	0	0	1
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	0

Truth Table



Binary Decision Tree



Binary Decision Diagram

Function Evaluation Examples:

$$f(0,1,0) = 0$$

$$f(1,1,0) = 1$$

Polyhedral Binary Decision Diagrams (PBDDs)

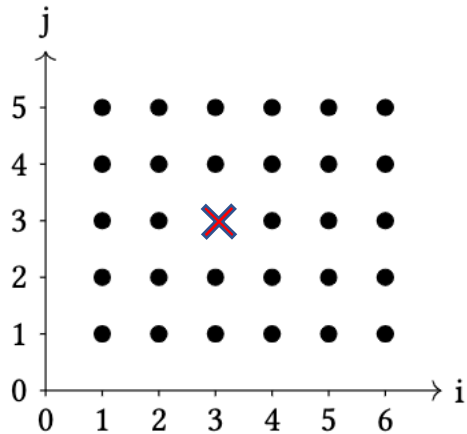
- ❖ In PBDDs, we consider the Boolean function being evaluated as the indicator of whether a point is in the polyhedral set.
- ❖ Variables correspond to constraints
- ❖ Point in iteration space corresponds to indicator vector of which constraints the point satisfies
- ❖ Function evaluation is same as in the case of BDDs

Polyhedral Binary Decision Diagrams (PBDDs)

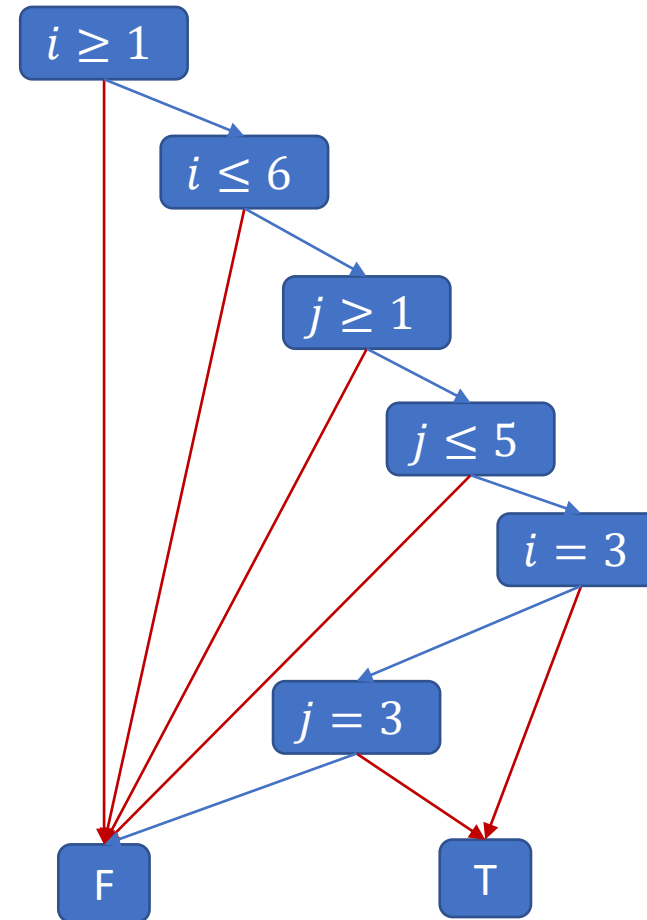
Program (recall)

```
for (int i = 1; i <= 6; ++i)
  for (int j = 1; j <= 5; ++j)
    if (i != 3 || j != 3)
      Stmt(i,j);
```

Statement Domain (recall)



PBDD Representation



Example Evaluations:

$(0,1) \rightarrow F$

$(2,4) \rightarrow T$

$(3,3) \rightarrow F$

Distinctions from BDDs and QuASTs

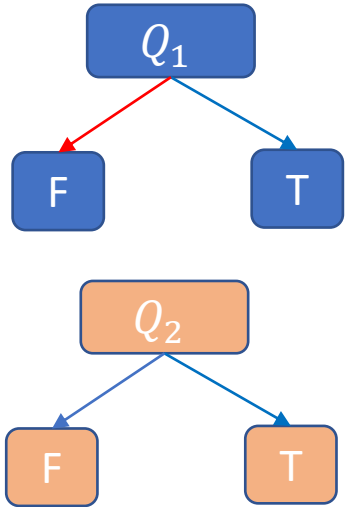
❖ PBDD vs BDD

- ❖ An assumption of BDDs is that all the input variables are independent
- ❖ For PBDDs that have not been simplified, conflicting constraints can happen
- ❖ We allow non-simplified PBDDs as simplification steps often computationally intensive

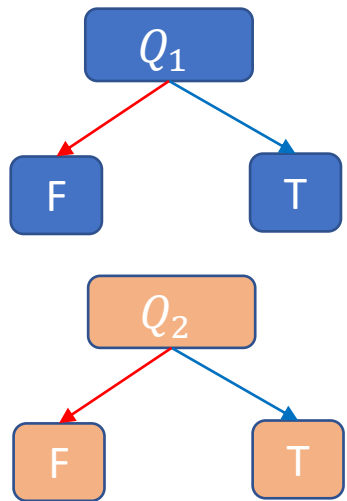
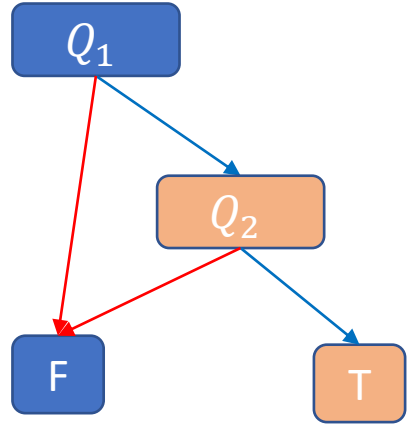
❖ PBDD vs QuAST

- ❖ Quasi-Affine Solution Tree
- ❖ Used to describe the piece-wise defined solution of the lexicographic minimum of a parametric Z-polyhedron.
- ❖ In contrast to PBDD, it is defined by a context-free grammar.
- ❖ QuASTs are always trees instead of DAGs.
- ❖ QuAST is vector-valued whereas PBDD is Boolean-valued (indicator function).

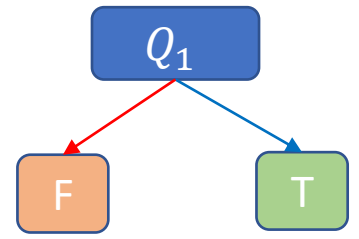
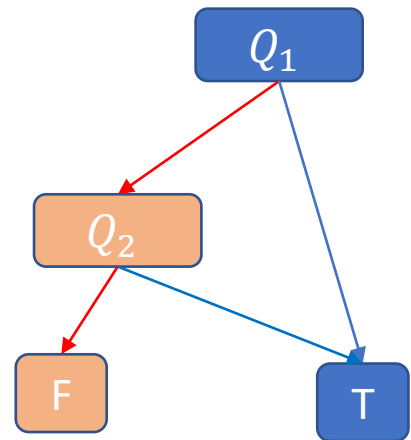
Simple Set Operations



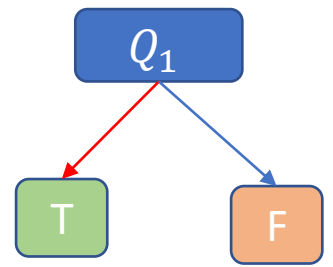
Intersect



Union



Complement



Can be implemented to be extremely efficient!!

Simple Set Operations

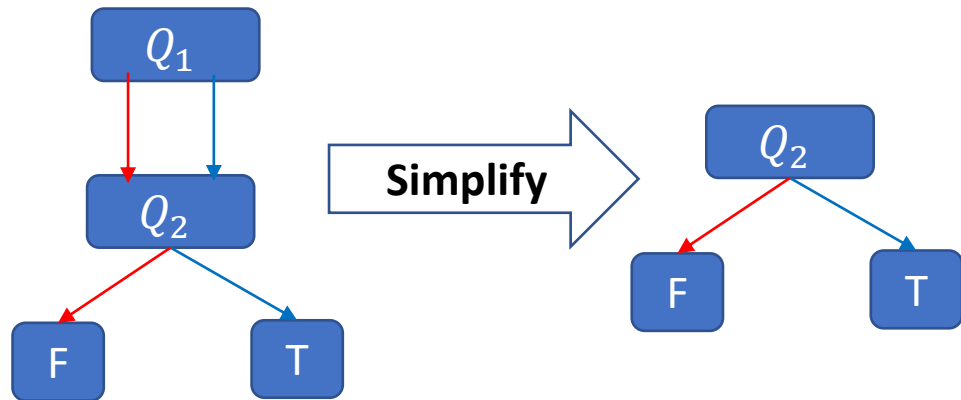
- ❖ Main paper also defines and gives algorithms for various set operations:
 - ❖ Emptiness check
 - ❖ Subset check
 - ❖ Subtraction
 - ❖ Project Out

- ❖ Recursion + Memoization

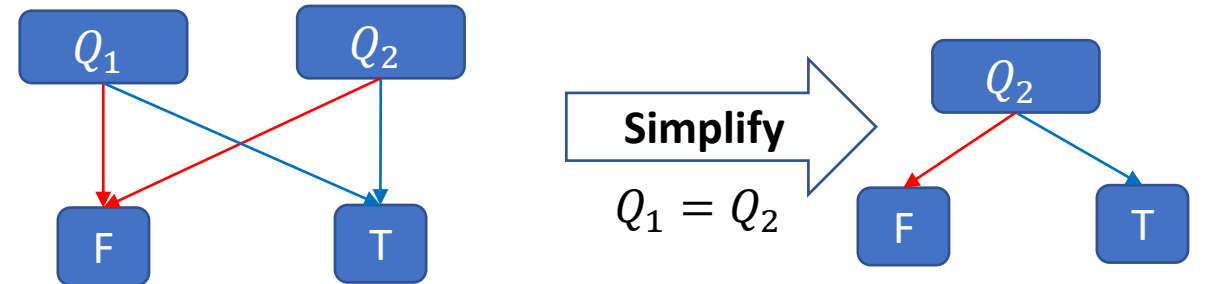
Note: Algorithms are recursive as our Python PoC implementation of a PBDD is inherently recursive.

Simplification Operations

- ❖ We use structural simplifications in order to control the sizes of PBDDs
- ❖ Conceptually simple but lead to big speedups
- ❖ Correctness follows from *Shannon Expansion* (see main paper)
- ❖ Experiments show that these are necessary to make representations tractable.



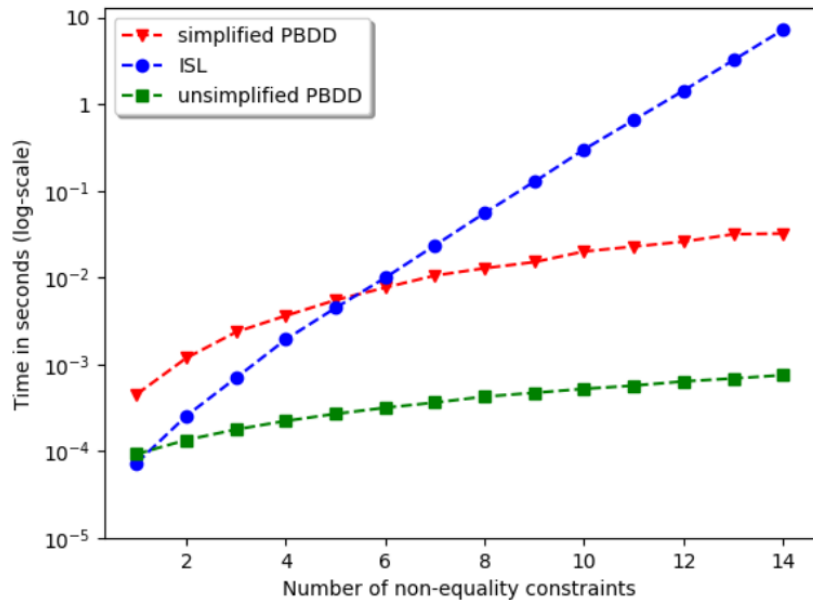
Redundant Nodes



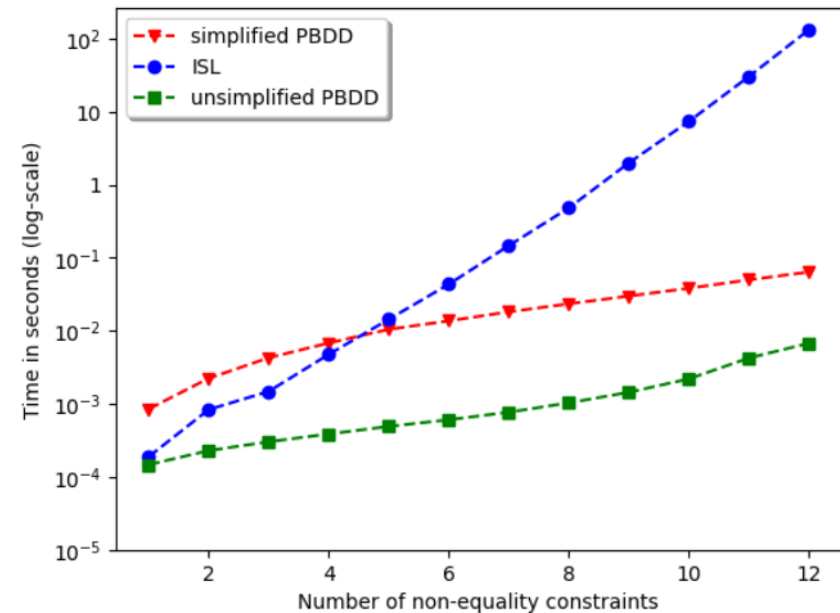
Isomorphic Nodes

Experiments I

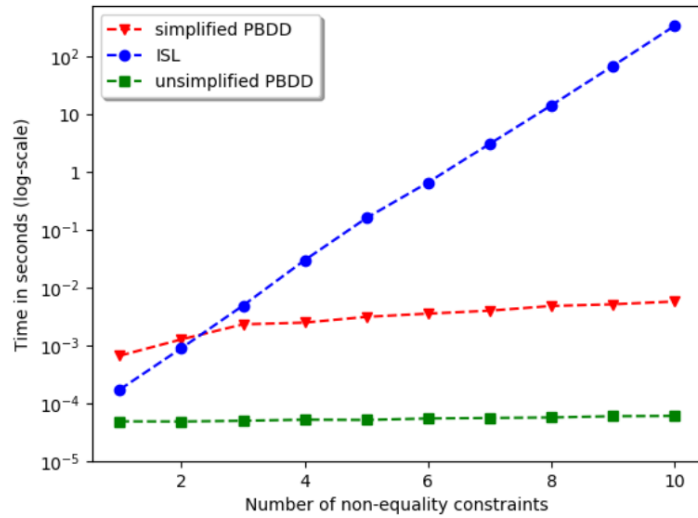
- (1) Input: k , constraints c_i
- (2) $P := \text{Universe}$
- (3) for $i = 1 \dots k$
 - (a) $P = \text{intersect}(P, \text{complement}(c_i))$
- (4) Output P



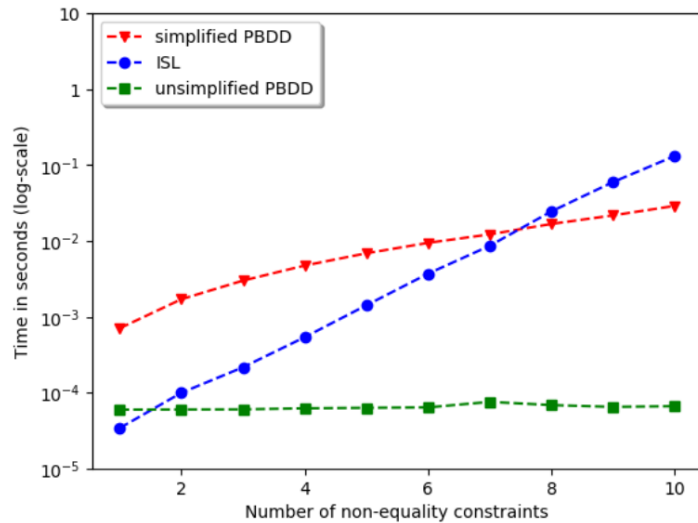
- (1) Input: k , constraints c_i
- (2) $P := \text{Universe}$
- (3) for $i = 1 \dots k$
 - (a) $P = \text{subtract}(P, \text{intersect}(c_i, P))$
- (4) Output P



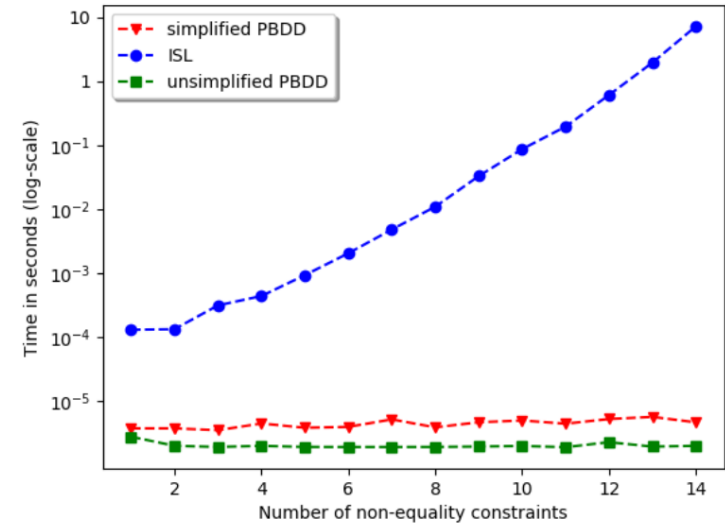
Experiments II



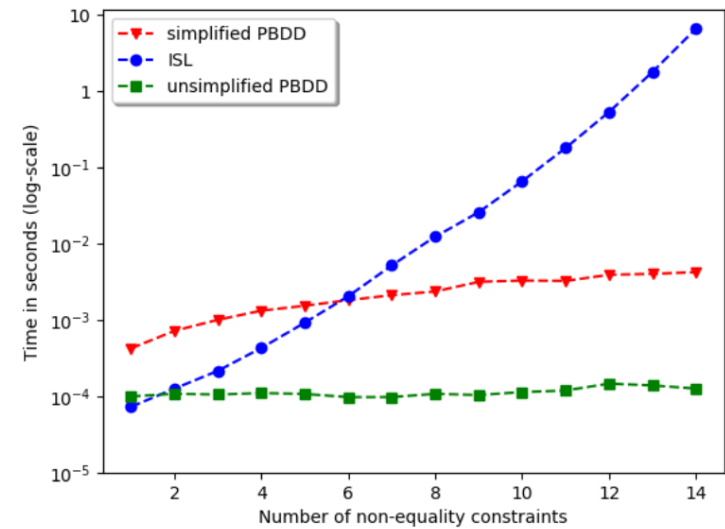
(a) Non-convex Intersection



(b) Non-convex Union

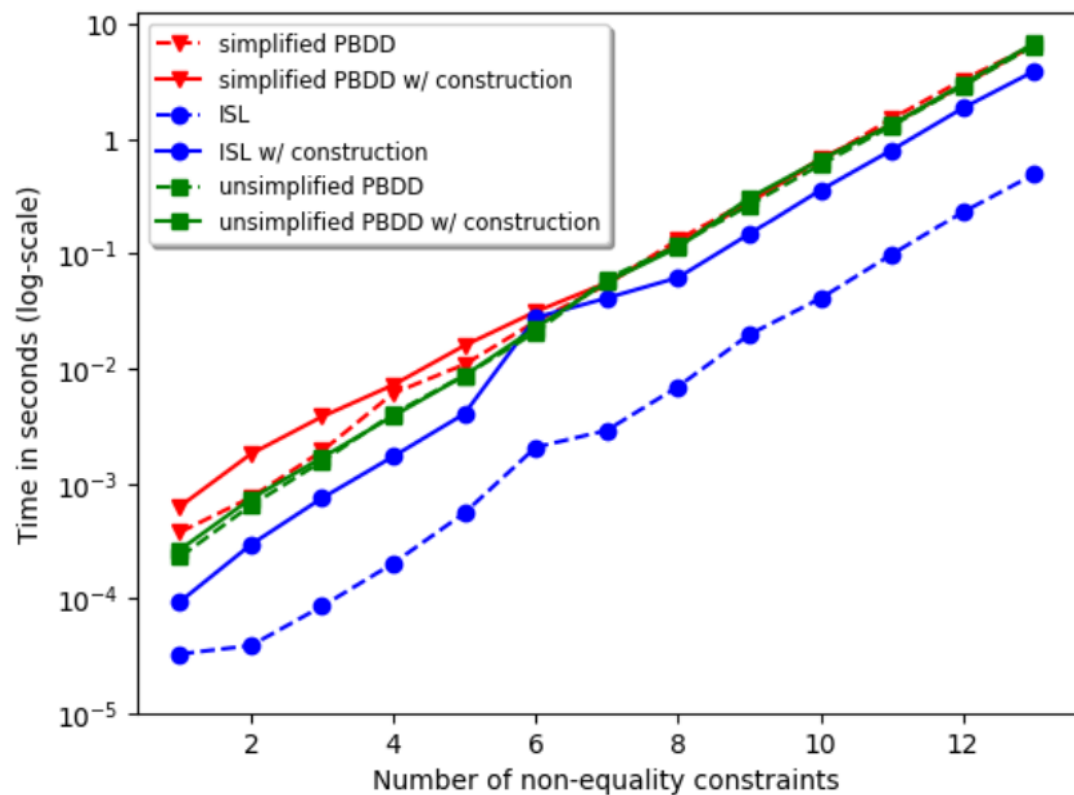


(c) Non-convex Complement

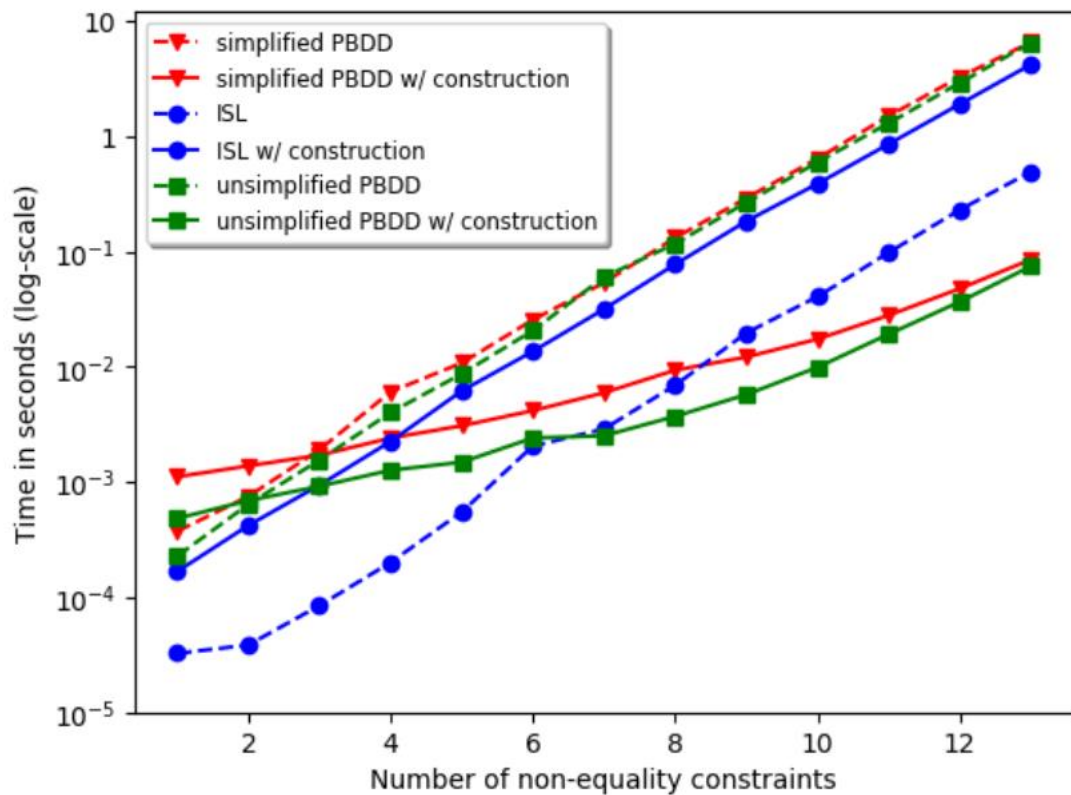


(d) Non-convex Subtraction

Experiments III



(a) Non-convex project-out



(b) Non-convex emptiness check

Future Work

❖ Implementation that does not rely on ISL intermediate data representation

❖ Parallelism

❖ Operational Efficiency

❖ Simplifications

❖ Decision Orderings

❖ Approximations

❖ Typed Roots

❖ Adapt PBDDs for ILP solvers

See main paper for discussions on each!